

Distributed Inference and Management of Future Cyber-Physical Networks

Georgios B. Giannakis

*Digital Technology Center and Dept. of ECE
University of Minnesota*

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Prof. N. Gatsis (UTSA), V. Kekatos (VaTech), H. Zhu (UIUC),
S. Dhople (UMN), Dr. E. Dall'Anese (NREL), and L. Zhang (UMN)

SMART GRID: Advanced infrastructure and information technologies (Cyber) to enhance the electrical power network (Physical)



controllable



resilient



efficient



participation



sustainable



self-restoring



green



situational awareness

Enabling technology advances



distributed generation
micro-grids



renewables

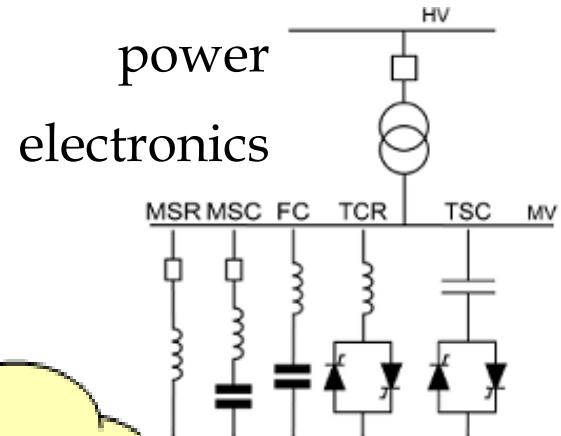
sensing/metering



demand response



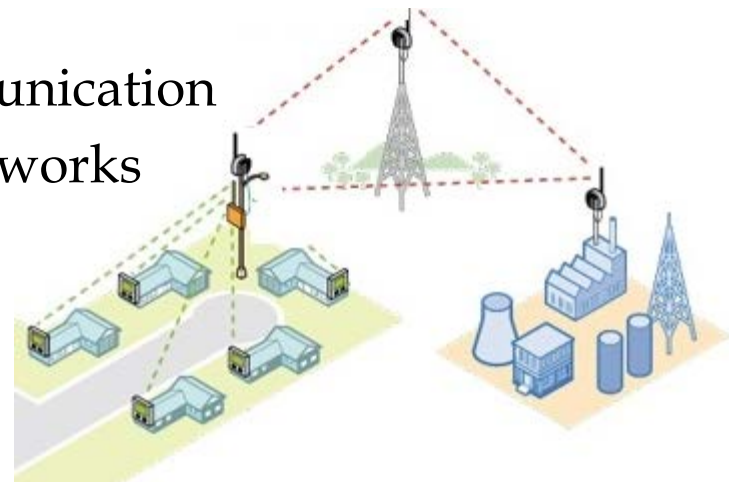
**Optimization, learning,
and signal processing
toolbox**



electric vehicles



communication
networks



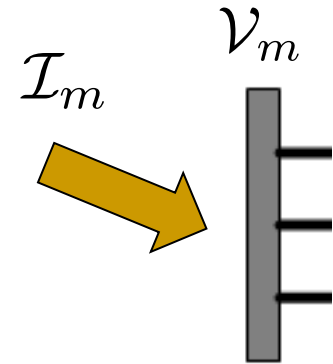
Outline

- Distributed and robust power system state estimation (PSSE)
- Distributed optimal power flow (OPF)
- Distributed demand response (DR)
- Distributed electric vehicle (EV) charging

Complex power

- Power injection to bus m

$$S_m = V_m I_m^* = P_m + jQ_m$$



- (Re) active power generated or consumed at a bus
- Power flow over line (m, n) $S_{mn} = V_m I_{mn}^* = P_{mn} + jQ_{mn}$
- Multivariate nodal power model (*quadratic* in \mathbf{v})

$$\mathbf{s} = \text{diag}(\mathbf{v}) \mathbf{i}^* = \text{diag}(\mathbf{v}) \mathbf{Y}^* \mathbf{v}^*$$

Diagram illustrating the multivariate nodal power model equation $\mathbf{s} = \text{diag}(\mathbf{v}) \mathbf{i}^* = \text{diag}(\mathbf{v}) \mathbf{Y}^* \mathbf{v}^*$. Blue arrows point from descriptive text to the variables in the equation:

- concatenating $\{V_m\}$ points to \mathbf{v} in $\text{diag}(\mathbf{v})$
- bus admittance matrix points to \mathbf{Y}^*
- concatenating $\{S_m\}$ points to \mathbf{s}
- concatenating $\{I_m\}$ points to \mathbf{i}^*



Power system state estimation

Motivation for PSSE

Goal: Given meter readings and grid parameters, find state vector \mathbf{v}

- Quantities of interest expressible as functions of bus voltages in \mathbf{v}
- PSSE is of paramount importance for
 - Situational awareness
 - Reliability analysis and planning
 - Load forecasting
 - Economic operations and billing
- Can be formulated as an estimation problem [Schweppe et al'70]

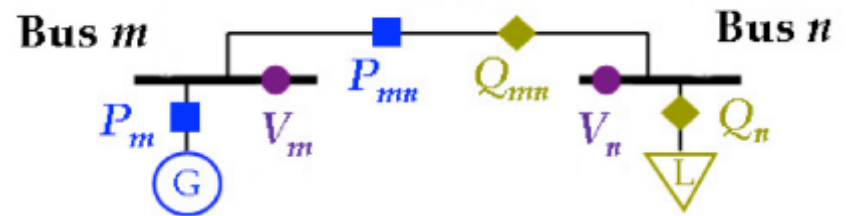
SCADA-based PSSE

- Supervisory control and data acquisition (SCADA) system
 - Terminals forward readings to control center (~4 secs)
 - Phases cannot be used due to timing mismatches

- Available measurements (M)

$$\{V_m, P_m, Q_m, P_{mn}, Q_{mn}, I_{mn}\}$$

$$\mathbf{z} = \mathbf{h}(\mathbf{v}) + \boldsymbol{\epsilon}$$



- Nonlinear (weighted) least-squares

$$\hat{\mathbf{v}} := \arg \min_{\mathbf{v}} \|\mathbf{z} - \mathbf{h}(\mathbf{v})\|^2$$

- Constraints

- Zero-injection buses $P_m = Q_m = 0$

- Feasible ranges $V_m^{\min} \leq V_m \leq V_m^{\max}$

Popular solvers

(M1) Gauss-Newton iterations

$$\hat{\mathbf{v}} := \arg \min_{\mathbf{v}} \|\mathbf{z} - \mathbf{h}(\mathbf{v})\|^2$$

- Approximate $\mathbf{h}(\mathbf{v}) \simeq \mathbf{h}(\mathbf{v}_k) + \mathbf{G}_k^T (\mathbf{v} - \mathbf{v}_k)$, \mathbf{G}_k : Jacobian at \mathbf{v}_k
- Linear LS in closed form $\mathbf{v}_{k+1} = \mathbf{v}_k + (\mathbf{G}_k \mathbf{G}_k^T)^{-1} \mathbf{G}_k (\mathbf{z} - \mathbf{h}(\mathbf{v}_k))$
- Cholesky factorization based remedies for numerical stability
- Sensitive to initialization; No convergence guarantee

(M2) Fast decoupled solver

- Active powers depend only on $\{\theta_m\}$; reactive only on $\{V_m\}$
- Approximate $(\mathbf{G}_k \mathbf{G}_k^T)^{-1}$ at *flat voltage profile* $\mathbf{v} = \mathbf{1} + j\mathbf{0}$

Semidefinite relaxation

- Rectangular coordinates: measurements are *quadratic* in \mathbf{v}

$$P_m + jQ_m = \mathcal{V}_m \mathcal{I}_m^* = \mathbf{e}_m^T \mathbf{v} (\mathbf{Y} \mathbf{v})^H \mathbf{e}_m = \text{Tr}(\underbrace{\mathbf{Y}^H \mathbf{e}_m \mathbf{e}_m^T}_{\mathbf{H}_m} \mathbf{v} \mathbf{v}^H)$$

- Yet *linear* in $\mathbf{V} = \mathbf{v} \mathbf{v}^H$

$$\min_{\mathbf{v}} \sum_{m=1}^M (z_m - h_m(\mathbf{v}))^2$$



$$\begin{aligned} & \min_{\mathbf{V}} \sum_{m=1}^M (z_m - \text{Tr}(\mathbf{H}_m \mathbf{V}))^2 \\ & \text{s.to } \mathbf{V} \succeq \mathbf{0} \text{ and } \text{rank}(\mathbf{V}) = 1 \end{aligned}$$

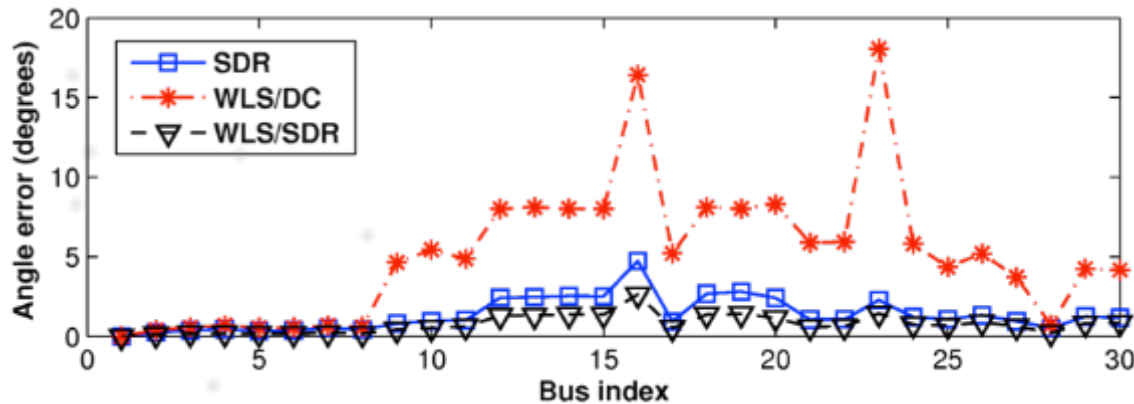
- SDR popular in SP and communications [Goemans et al'95]
- SDR for SE [Zhu-GG'11], SDR for OPF [Bai et al'08], [Lavaei-Low'11]
 - Generalizations include PMU data, and robust SDR-based PSSE
 - (Near-)optimal regardless of initialization; polynomial complexity $O(N^{4.5} \log(1/\epsilon))$

Numerical tests

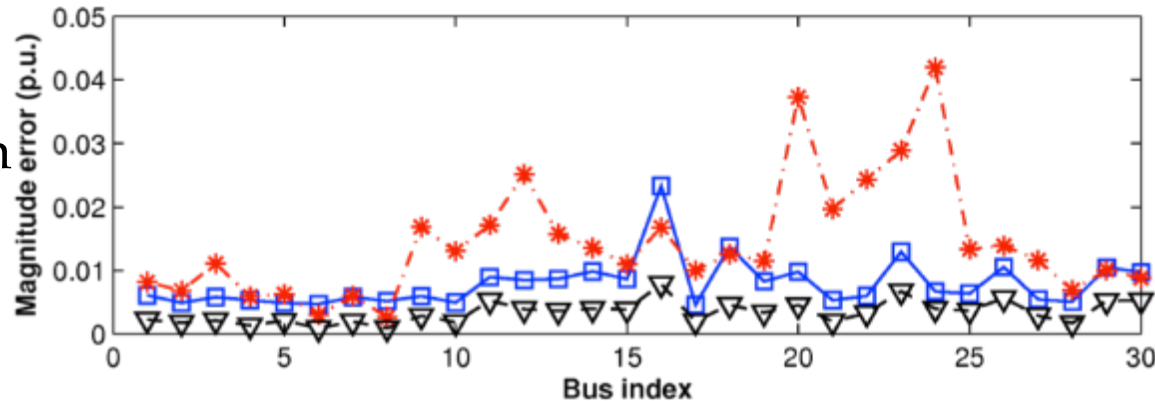
- IEEE 30, 57, and 118-bus benchmarks
- $V_m \sim \mathcal{N}(1, 0.01)$, $\theta_m \sim \mathcal{U}[-\theta, \theta]$

Average running time in secs.

# of buses	WLS	SDR
30	0.216	1.62
57	0.558	4.32
118	2.87	21.6

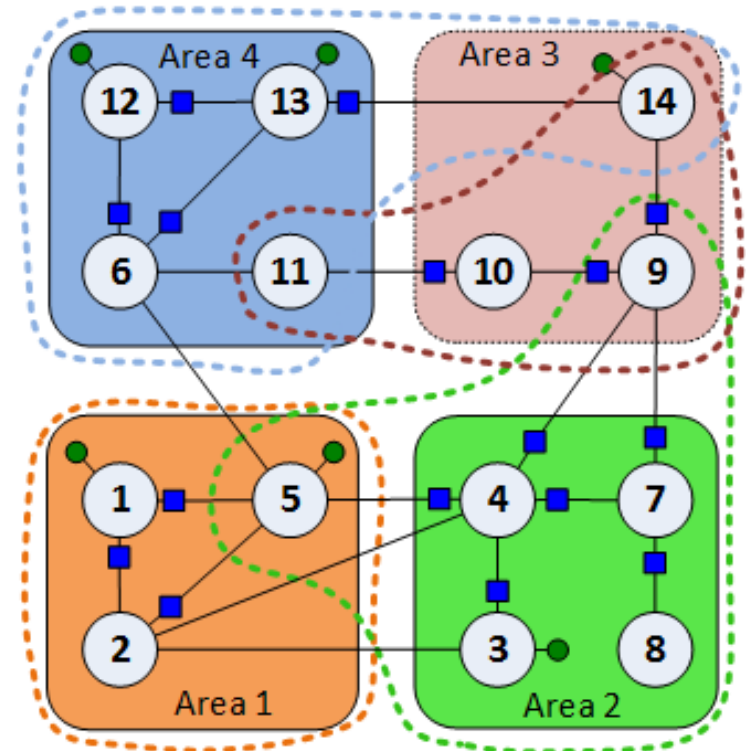


- Closer to global optimum at higher complexity



Decentralized PSSE - motivation

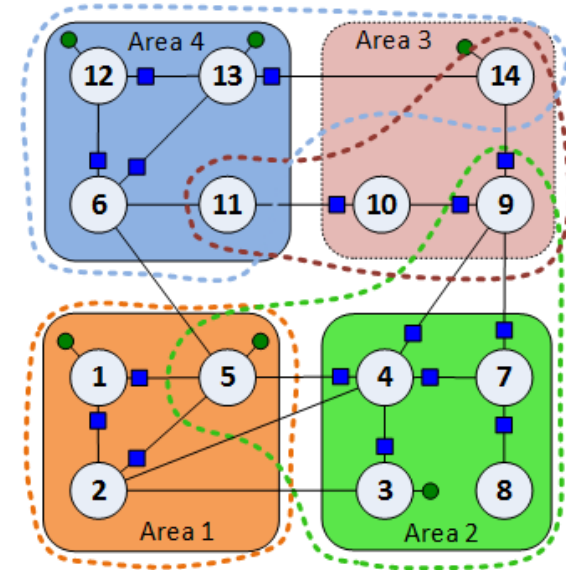
- Scalable with control area size, and privacy preserving
- Area 2 buses (states): $\{3,4,7,8\}$
- Area 2 collects flow measurements $\{(4,5), (4,9), (7,9)\}...$
- *Option 1*: Ignore tie-line meters
 - statistically suboptimal
 - observability at risk (bus 11)
 - tie-line mismatches (trading)
- *Option 2*: Augment \mathbf{v}_2 to $\{3,4,7,8,5,9\}$
 - consent with neighbors on shared states



Cost decomposition

- Include tie-line buses to split local LS cost per $\mathcal{N}_{(k)}$

$$f_k(\mathbf{V}_{(k)}) := \sum_{\ell=1}^{M_k} \left[z_k^\ell - \text{Tr}(\mathbf{H}_{(k)}^\ell \mathbf{V}_{(k)}) \right]^2$$



$$\mathcal{N}_{(2)} := \mathcal{N}_2 \cup \{5, 9\}$$

$$\hat{\mathbf{V}} := \arg \min_{\mathbf{V}} \sum_{k=1}^K \sum_{\ell=1}^{M_k} \left[z_k^\ell - \text{Tr}(\mathbf{H}_k^\ell \mathbf{V}) \right]^2$$

s.to $\mathbf{V} \succeq \mathbf{0}$

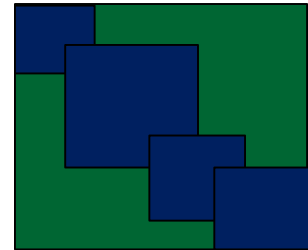


$$\hat{\mathbf{V}} := \arg \min_{\mathbf{V}} \sum_{k=1}^K f_k(\mathbf{V}_{(k)})$$

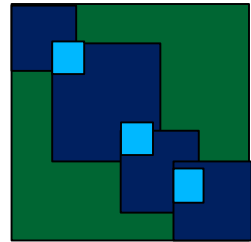
s.to $\mathbf{V} \succeq \mathbf{0}$

Challenge: as $\{\mathcal{N}_{(k)}\}$ overlap partially, PSD constraint couples $\{\mathbf{V}_{(k)}\}$

Blessing: overlap \rightarrow global; no overlap: $\mathbf{V} \succeq \mathbf{0} \Leftrightarrow \mathbf{V}_{(k)} \succeq \mathbf{0}, \forall k$



Distributed SDP for PSSE



- If *graph* with *areas-as-nodes* and *overlaps-as-edges* is a *tree*, then

$$\hat{\mathbf{V}} := \arg \min_{\mathbf{V}} \sum_{k=1}^K f_k(\mathbf{V}_{(k)})$$

(C-SDP)

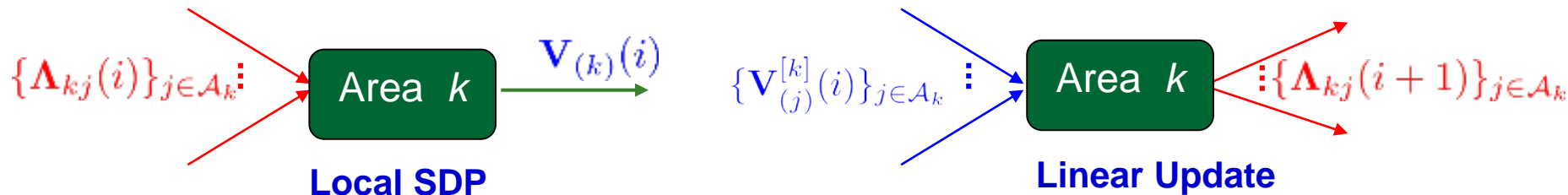
s.to $\mathbf{V} \succeq \mathbf{0}$



$$\{\hat{\mathbf{V}}_{(k)}\} := \arg \min_{\{\mathbf{V}_{(k)}\}} \sum_{k=1}^K f_k(\mathbf{V}_{(k)})$$

s.to $\mathbf{V}_{(k)} \succeq \mathbf{0}, \mathbf{V}_{(k)}^{[j]} = \mathbf{V}_{(j)}^{[k]}$

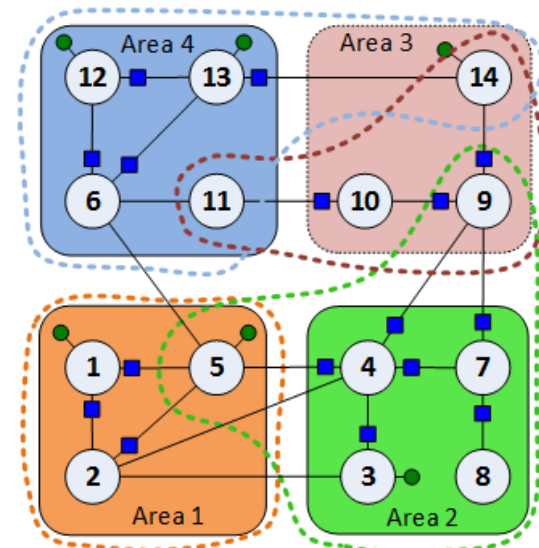
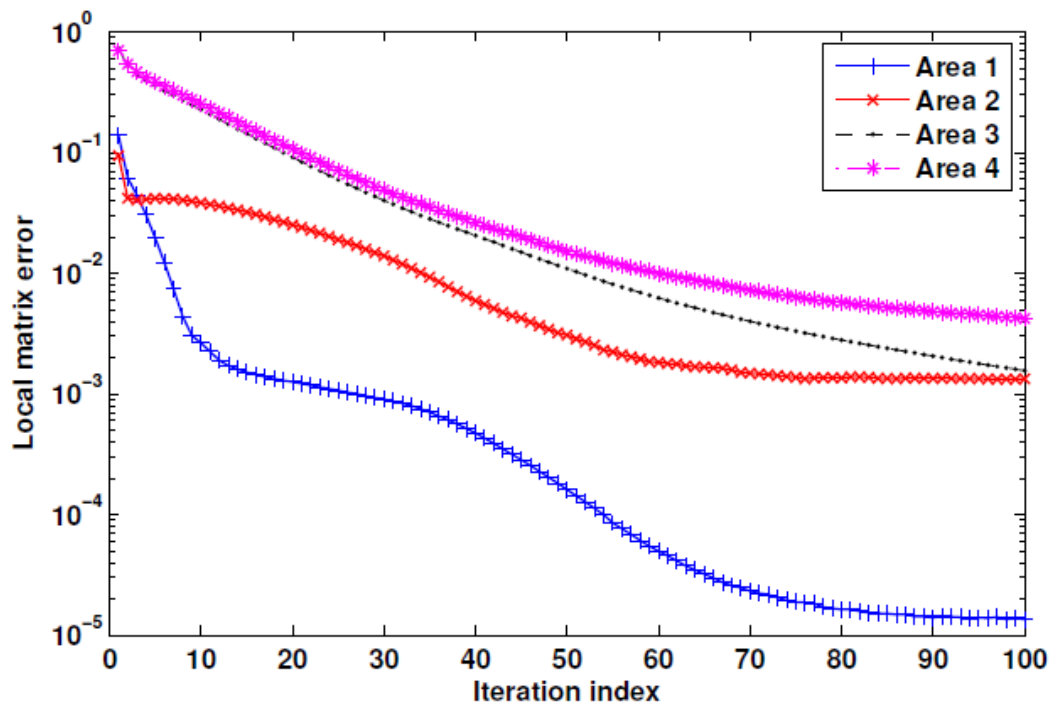
- ADMM [Glowinski-Marrocco'75]; for D-Estimation [Schizas-Giannakis'07]
 - Iterates between local variables and multipliers per equality constraint



- Converges $\mathbf{V}_{(k)}(i) \rightarrow \hat{\mathbf{V}}_{(k)}$ even for noisy-async. links [Schizas-GG'08], [Zhu-GG'09]

ADMM convergence in action

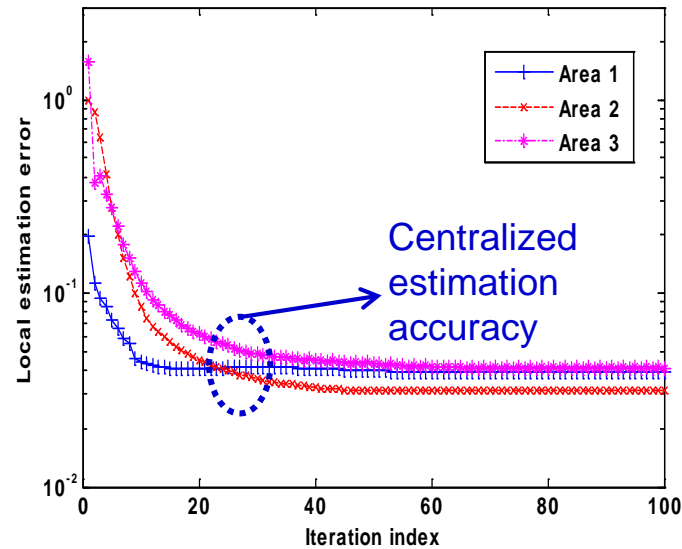
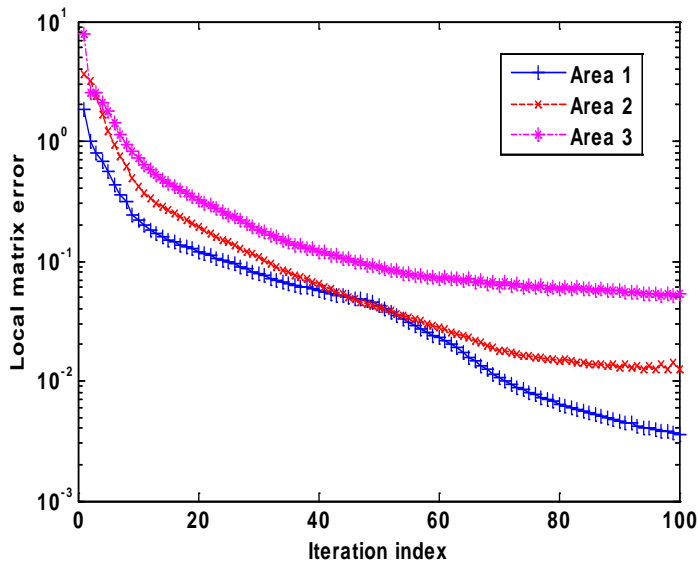
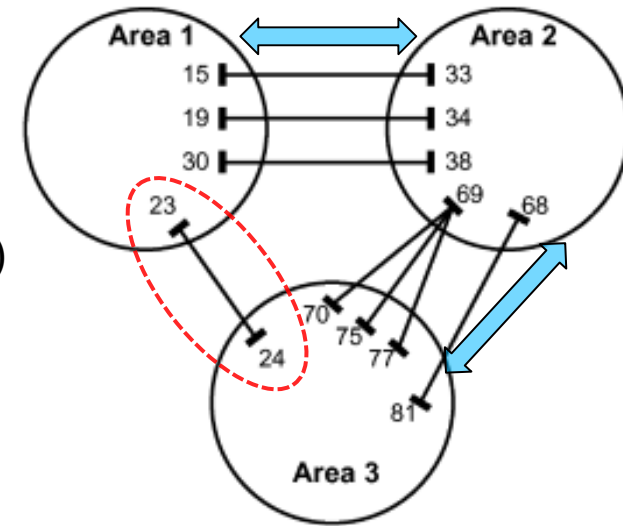
- IEEE 14-bus grid with 4 areas; 5 meters on tie-lines



- Errors $\|\mathbf{V}_{(k)}(i) - \hat{\mathbf{V}}_{(k)}\|_F$ vanish asymptotically

118-bus test case

- Triangular configuration [Min-Abur'06]
 - Power flow meters on all tie lines except for (23, 24)
- ➡ graph of areas is a tree



- Local norms

$$\|\mathbf{v}_{(k)}(i) - \mathbf{v}_{(k)}\|_2$$

converge in only
20 iterations!

Decentralized PSSE for linear models

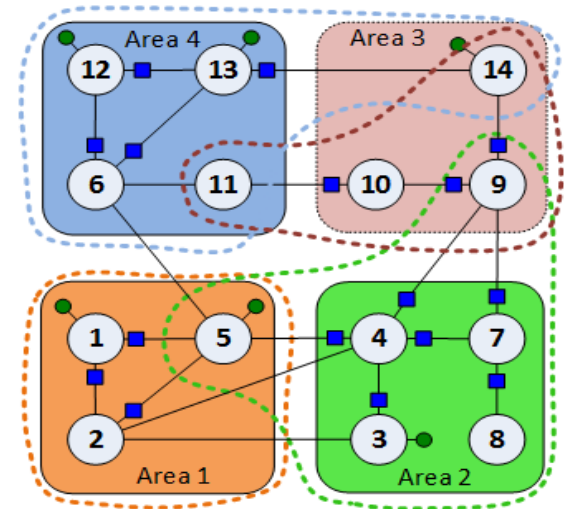
- Local linear(ized) model $\mathbf{z}_k = \mathbf{H}_k \mathbf{x}_k + \mathbf{n}_k$

- Regional PSSEs

$$\min_{\mathbf{x}_k \in \mathcal{X}_k} f_k(\mathbf{x}_k)$$

- Coupled local problems

$$\begin{aligned} \min_{\{\mathbf{x}_k\}} & \sum_{k=1}^K f_k(\mathbf{x}_k) \\ \text{s.t.} & \mathbf{x}_k[l] = \mathbf{x}_l[k] \end{aligned}$$



$$\mathcal{S}_2 := \mathcal{N}_{(2)} \setminus \mathcal{N}_2$$

$$\mathbf{S1.} \quad \mathbf{x}_k^{t+1} \Leftarrow \arg \min_{\mathbf{x}_k \in \mathcal{X}_k} f_k(\mathbf{x}_k) + \frac{c}{2} \sum_{i \in \mathcal{S}_k} (x_k(i) - p_k^t(i))^2$$

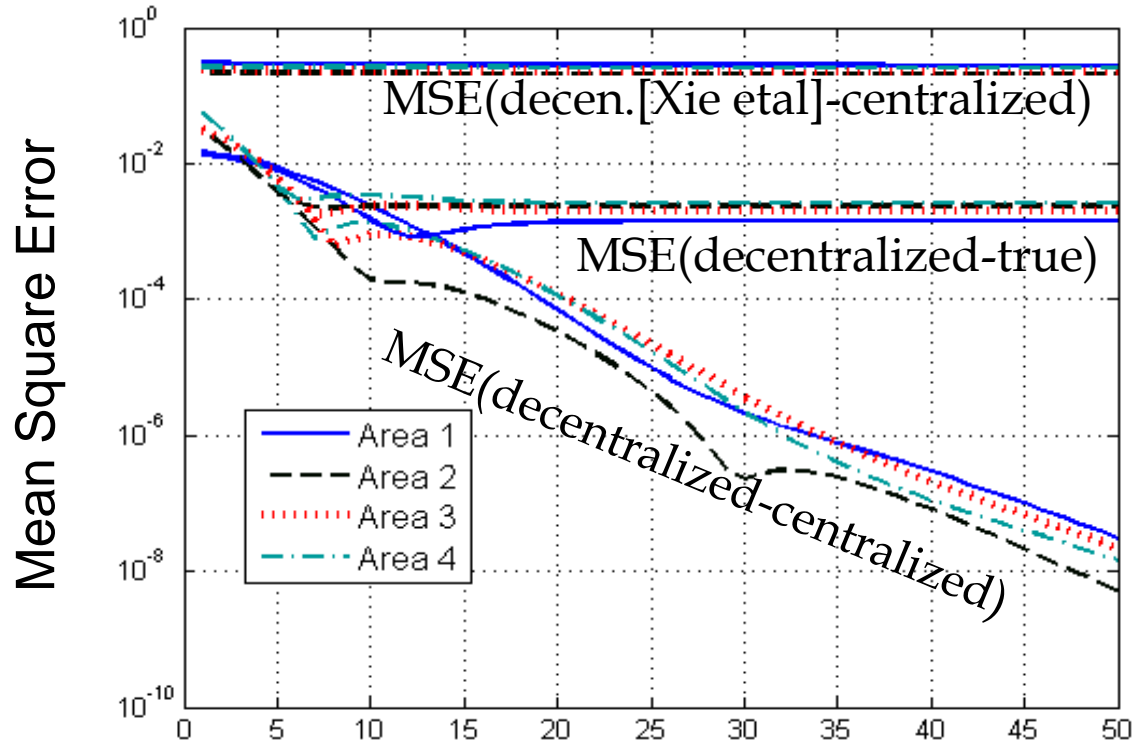
$$\mathbf{S2.} \quad p_k^{t+1}(i) \Leftarrow p_k^t(i) + \left(x_l^{t+1}[i] - \frac{x_k^t(i) + x_l^t[i]}{2} \right)$$

- ADMM solver: convergent with minimal exchanges and privacy-preserving

Simulated test

S1. $\mathbf{x}_k^{t+1} \Leftarrow (\mathbf{H}_k^T \mathbf{H}_k + c \cdot \mathbf{D}_k)^{-1} (\mathbf{H}_k^T \mathbf{z}_k + c \cdot \mathbf{D}_k \mathbf{p}_k^t)$ $[\mathbf{D}_k]_{ii} = |\mathcal{S}_k^i|$

S2. $p_k^{t+1}(i) \Leftarrow p_k^t(i) + \left(x_l^{t+1}[i] - \frac{x_k^t(i) + x_l^t[i]}{2} \right)$



Decentralized bad data cleansing

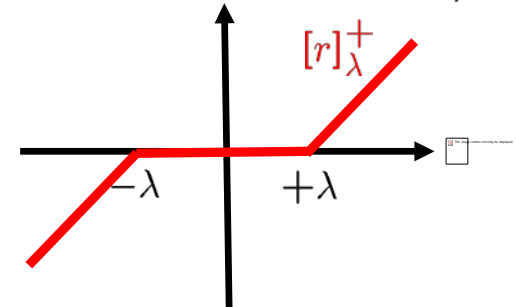
$$\mathbf{z} = \mathbf{H}\mathbf{x} + \mathbf{n} + \mathbf{o}$$

- Reveal *single* and *block* outliers via

$$\begin{aligned} f(\mathbf{x}) &:= \min_{\mathbf{o}} \frac{1}{2} \|\mathbf{z} - \mathbf{H}\mathbf{x} - \mathbf{o}\|_2^2 + \lambda \|\mathbf{o}\|_1 \\ &= \sum_{m=1}^M h(z_m - \mathbf{h}_m^T \mathbf{x}) \end{aligned}$$

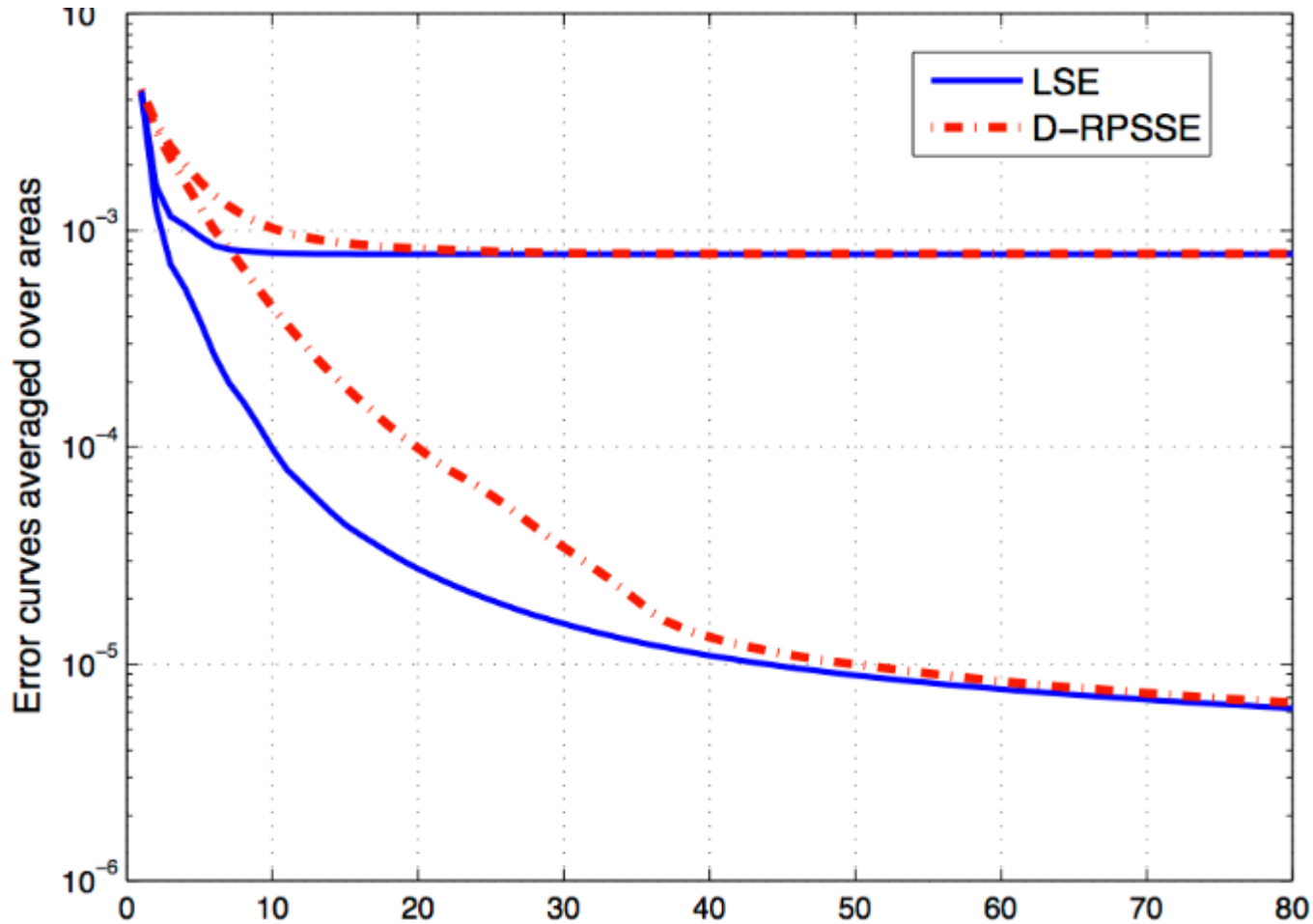
$$\mathbf{S1.} \quad \mathbf{x}_k^{t+1} \Leftarrow (\mathbf{H}_k^T \mathbf{H}_k + c \cdot \mathbf{D}_k)^{-1} (\mathbf{H}_k^T (\mathbf{z}_k - \mathbf{o}_k^t) + c \cdot \mathbf{D}_k \mathbf{p}_k^t)$$

$$\mathbf{S2.} \quad \mathbf{o}_k^{t+1} \Leftarrow [\mathbf{z}_k - \mathbf{H}_k \mathbf{x}_k^{t+1}]_{\lambda}^+$$



$$\mathbf{S3.} \quad p_k^{t+1}(i) \Leftarrow p_k^t(i) + \left(x_l^{t+1}[i] - \frac{x_k^t(i) + x_l^t[i]}{2} \right)$$

D-PSSE on a 4,200-bus grid





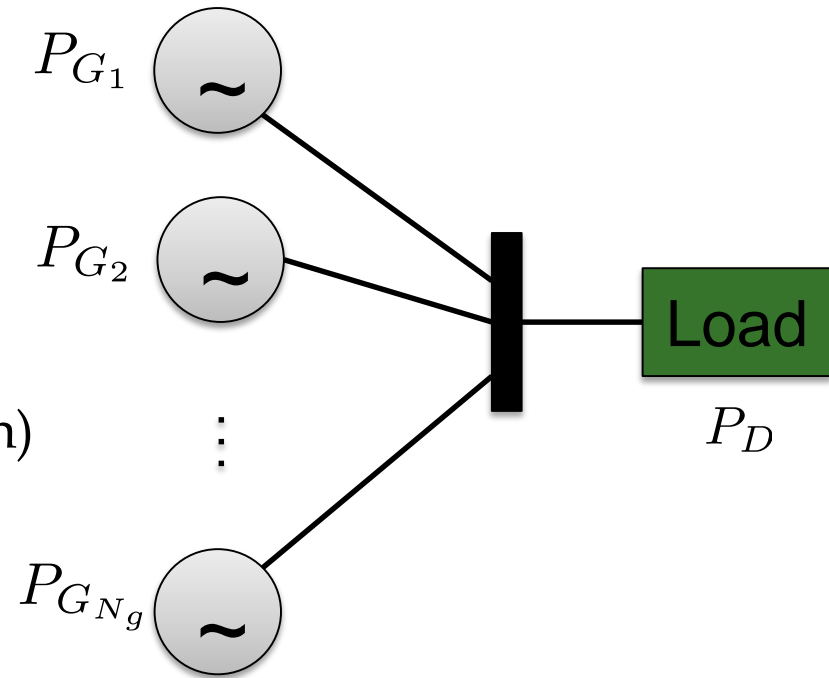
Optimal Power Flow

A. Gomez-Exposito, A. J. Conejo, and C. Canizares, *Electric Energy Systems: Analysis and Operation*, CRC, 2009.

A. J. Wood and B. F. Wollenberg, *Power Generation, Operation, and Control*, 2nd ed., Wiley, 1996.

Generation cost

- Thermal generators
- Power output P_{G_i} (MW)
- Generation cost $C_i(P_{G_i})$ (\$/h or €/h)



Economic dispatch (ED): Find most economically generated power output to serve given load

- ED typically solved every 5-10 minutes

AC optimal power flow

- **Motivation:** Minimize generation cost respecting physical constraints


$$\min_{\mathbf{p}_G, \mathbf{q}_G, \mathbf{v}} \sum_{m=1}^{N_b} C_m(P_{G_m})$$

$$\text{subj. to } \mathbf{p}_G - \mathbf{p}_D + j(\mathbf{q}_G - \mathbf{q}_D) = \text{diag}(\mathbf{v})(\mathbf{Y}\mathbf{v})^*$$

$$|\text{Re}\{\mathcal{S}_{mn}\}| \leq f_{mn}^{\max}; |\mathcal{S}_{mn}| \leq S_{mn}^{\max}$$

$$V_m^{\min} \leq |\mathcal{V}_m| \leq V_m^{\max}$$

$$\mathbf{p}_G^{\min} \leq \mathbf{p}_G \leq \mathbf{p}_G^{\max}; \mathbf{q}_G^{\min} \leq \mathbf{q}_G \leq \mathbf{q}_G^{\max}$$

- Quadratic equality constraints  nonconvex problem
- Traditional approaches rely on KKT conditions

SDP relaxation

$$\mathcal{V}_m \mathcal{I}_m^* = \mathbf{e}_m^H \mathbf{v} (\mathbf{Y} \mathbf{v})^H \mathbf{e}_m = \text{tr}[\mathbf{e}_m^H \mathbf{v} \mathbf{v}^H \mathbf{Y}^H \mathbf{e}_m] = \text{tr}[\mathbf{Y}^H \mathbf{e}_m \mathbf{e}_m^H \mathbf{v} \mathbf{v}^H]$$

- Nodal balance constraint linear in $\mathbf{V} := \mathbf{v} \mathbf{v}^H$

$$P_{G_m} - P_{D_m} + j(Q_{G_m} - Q_{D_m}) = \text{tr}[\mathbf{Y}^H \mathbf{e}_m \mathbf{e}_m^H \mathbf{V}]$$

- Line flow and bus voltage constraints also linear in \mathbf{V}
- AC-OPF with variables $\mathbf{p}_G, \mathbf{q}_G, \mathbf{V}$ and additional constraints

$$\mathbf{V} \succeq \mathbf{0}$$

$$\text{rank}[\mathbf{V}] = 1$$

⇒ Nonconvex ⇒ Drop

- Works in many practical OPF instances and IEEE benchmarks
- Optimal in tree graphs [Lam et al'12]

AC OPF for multi-phase

- Power and voltage magnitude as linear functions of $\mathbf{V} := \mathbf{v}\mathbf{v}^H$
- Regulating constraints per node *and* per phase (can be unbalanced)

$$\begin{aligned} & \min_{\mathbf{V}, \mathbf{p}_G, \mathbf{q}_G} C(\mathbf{p}_G, \mathbf{q}_G) \\ & \text{sub. to } \{P_{G,n}^\phi, Q_{G,n}^\phi\} \in \mathcal{B}_n^\phi, \text{ and } \forall \phi \text{ } n \\ & \text{Tr}(\mathbf{\Phi}_n^\phi \mathbf{V}) = P_{G,n}^\phi - P_{\ell,n}^\phi \\ & \text{Tr}(\mathbf{\Psi}_n^\phi \mathbf{V}) = Q_{G,n}^\phi - Q_{\ell,n}^\phi \\ & (V^{\min})^2 \leq \text{Tr}(\mathbf{M}_n^\phi \mathbf{V}) \leq (V^{\max})^2 \\ & \mathbf{V} \succeq \mathbf{0}, \text{ rank}(\mathbf{V}) = 1 \end{aligned}$$

- $\text{rank}(\mathbf{V}^{\text{opt}}) = 1 \Rightarrow$ globally optimal AC OPF solution!

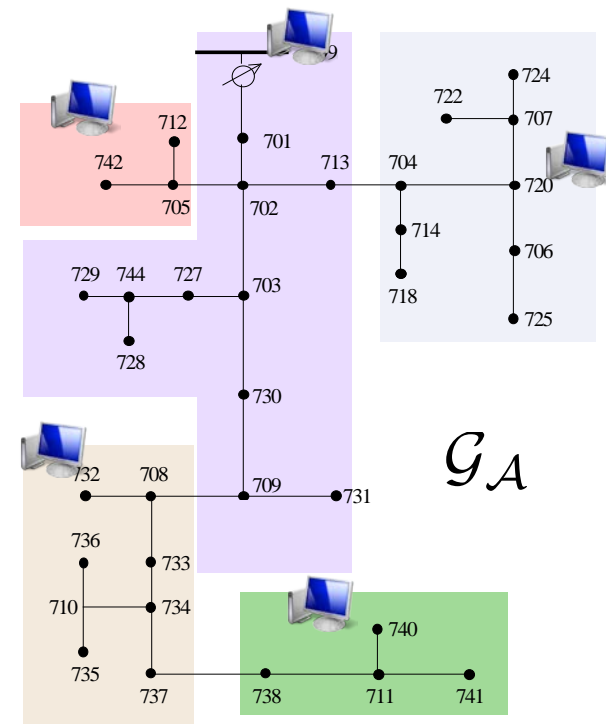
Distributed three-phase OPF

- Multi-area based on non-convex OPF [Kim-Baldick'97, Hug-Andersson'09, Erseghe'14]
- Node-to-node for single-phase systems [Zhang et al'12]
- Distributed SDP for three-phase systems

$$\begin{aligned} \min_{\mathbf{V}, \{\mathbf{s}_n\}} \quad & \sum_a C_a(\mathbf{V}^{(a)}, \mathbf{s}^{(a)}) \\ \text{sub. to} \quad & \{\mathbf{V}^{(a)}, \mathbf{s}^{(a)}\} \in \mathcal{B}^{(a)} \quad \forall a \end{aligned}$$

$$\mathbf{V} \succeq \mathbf{0}$$

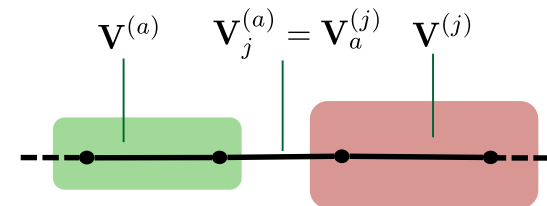
Challenge: PSD constraint couples local quantities! $\mathbf{Q}: \mathbf{V} \succeq \mathbf{0} \overset{?}{\longleftrightarrow} \{\mathbf{V}^{(a)} \succeq \mathbf{0}\}$



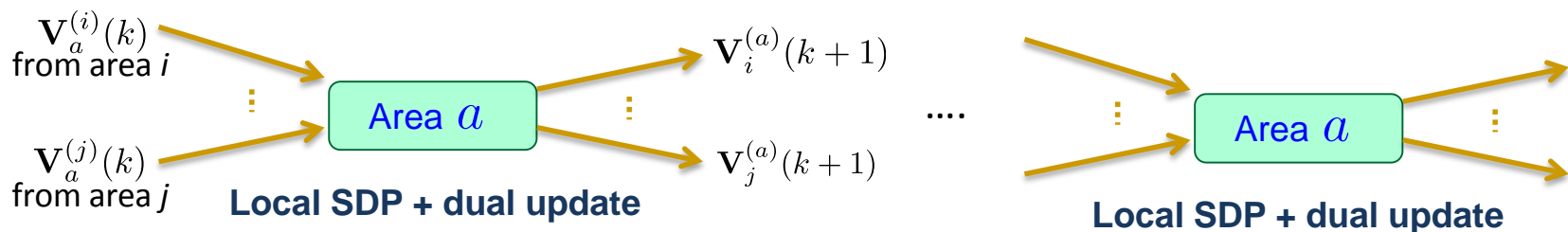
Topology-based decoupling

Result: If the *graph of areas* is a tree, without “loops” across areas, then the centralized PSD constraint *decouples*

$$\begin{aligned} \{\mathbf{V}_{\text{opt}}^{(a)}, \mathbf{s}_{\text{opt}}^{(a)}\} &= \arg \min_{\{\mathbf{V}^{(a)}\}, \{\mathbf{s}^{(a)}\}} \sum_a C_a(\mathbf{V}^{(a)}, \mathbf{s}^{(a)}) \\ \text{sub. to } &\{\mathbf{V}^{(a)}, \mathbf{s}^{(a)}\} \in \mathcal{B}^{(a)} \quad \forall a \\ &\mathbf{V}^{(a)} \succeq \mathbf{0} \quad \forall a \\ &\mathbf{V}_j^{(a)} = \mathbf{V}_a^{(j)} \quad \forall (a, j) \in \mathcal{E}_{\mathcal{A}} \end{aligned}$$

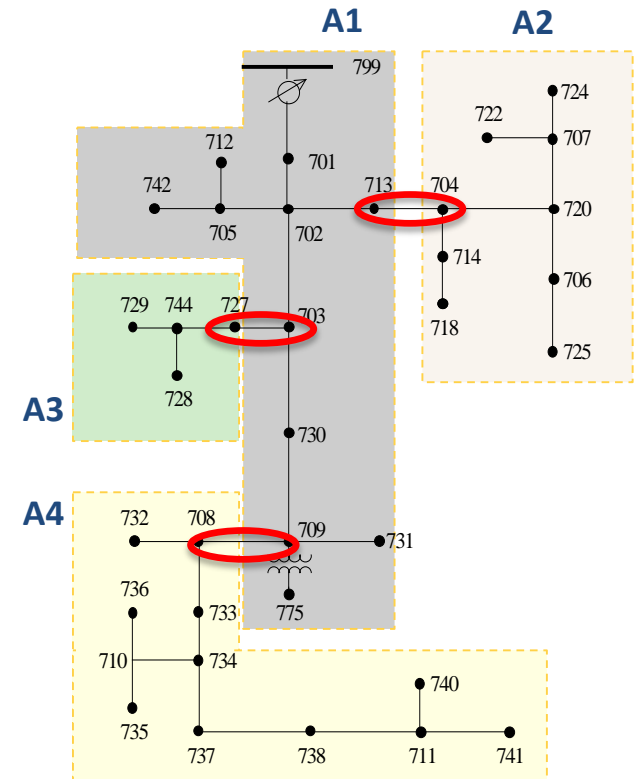
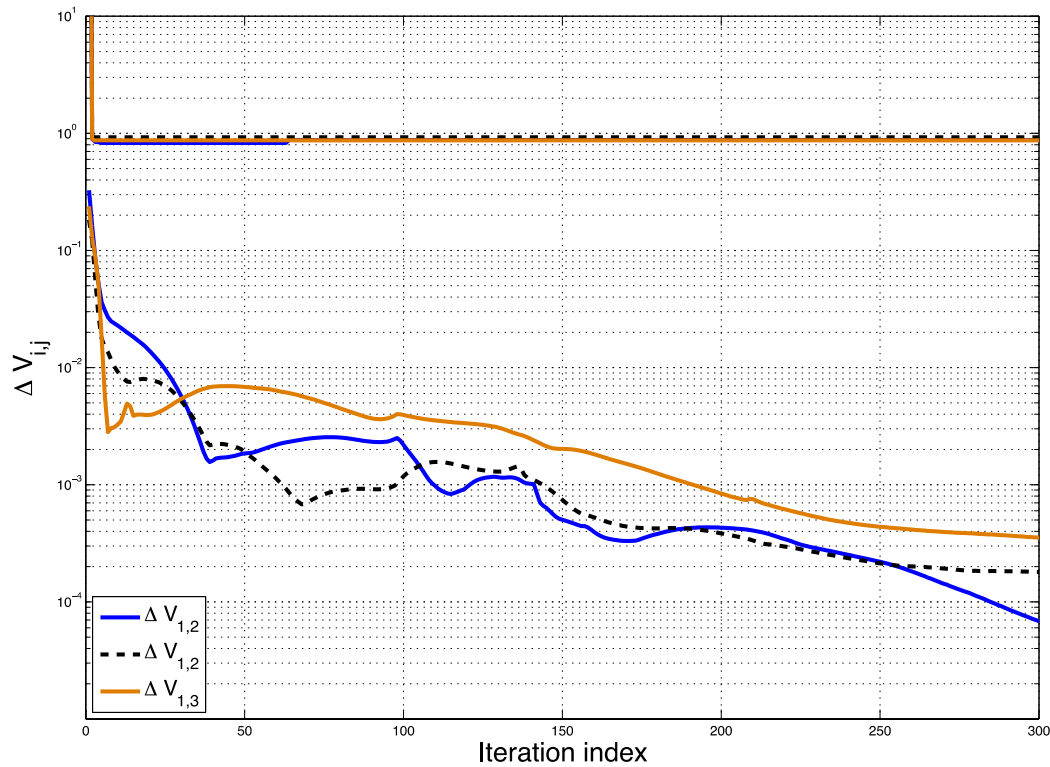


- Distributed solution via alternating direction method of multipliers



Illustrative test case

■ Consensus error



- Comparison with sub-gradient [Zhang et al'12]
- Convergence rate **does not** depend on area size



Demand Response

Motivation for DR

- Changes in electricity consumption by end-users in response to
 - Changes in electricity prices over time
 - Incentive payments at times of high wholesale prices

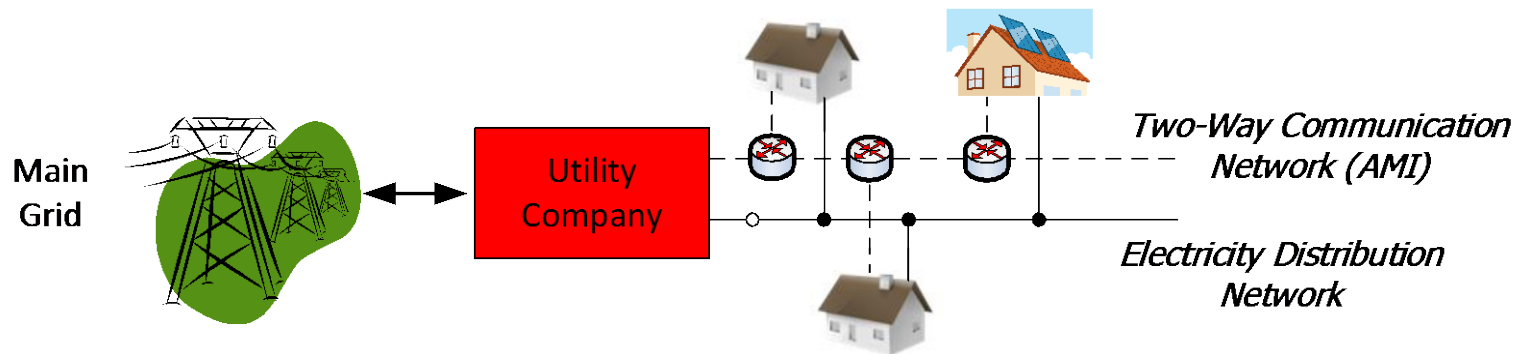
- Benefits of DR
 - Reduced demand reduces the potential of forced outages
 - Lower demand holds down electricity prices in spot markets
 - Can reduce the amount of generation and transmission assets

- DR programs
 - Incentive-based programs
 - Price-driven programs

Cooperative DR

- Set of users (residences) $\{1, \dots, R\}$ served by the same utility
- Set of smart appliances \mathcal{A}_r per user r
- Power consumption p_{ra}^t
- End-user utility function $U_{ra}(\mathbf{p}_{ra})$

- Cost of power procurement for utility company $C^t \left(\sum_{r=1}^R \sum_{a \in \mathcal{A}_r} p_{ra}^t \right)$



Social welfare maximization

$$\min_{\{\mathbf{p}_{ra}\}} \sum_{t=1}^T C^t \left(\sum_{r=1}^R \sum_{a \in \mathcal{A}_r} p_{ra}^t \right) - \sum_{r=1}^R \sum_{a \in \mathcal{A}_r} U_{ra}(\mathbf{p}_{ra})$$

subj. to $\mathbf{p}_{ra} \in \mathcal{P}_{ra}, a \in \mathcal{A}_r, r = 1 \dots, R$

- **Motivation:** Reduce peak demand respecting preferences of users
- Convexity depends on \mathcal{P}_{ra}
- Challenges
 - Scalable scheduling over AMI; and privacy issues

Solution approaches

- Gradient projection, block coordinate descent, dual decomposition
[Chen etal'12], [Mohsenian-Rad etal'10], [Papavasiliou etal'10],
[Samadi etal'11], [Gatsis-GG'12]

- **Dual decomposition:** Introduce variable s^t for total supplied power

$$\sum_{r=1}^R \sum_{a \in \mathcal{A}_r} p_{ra}^t \leq s^t$$

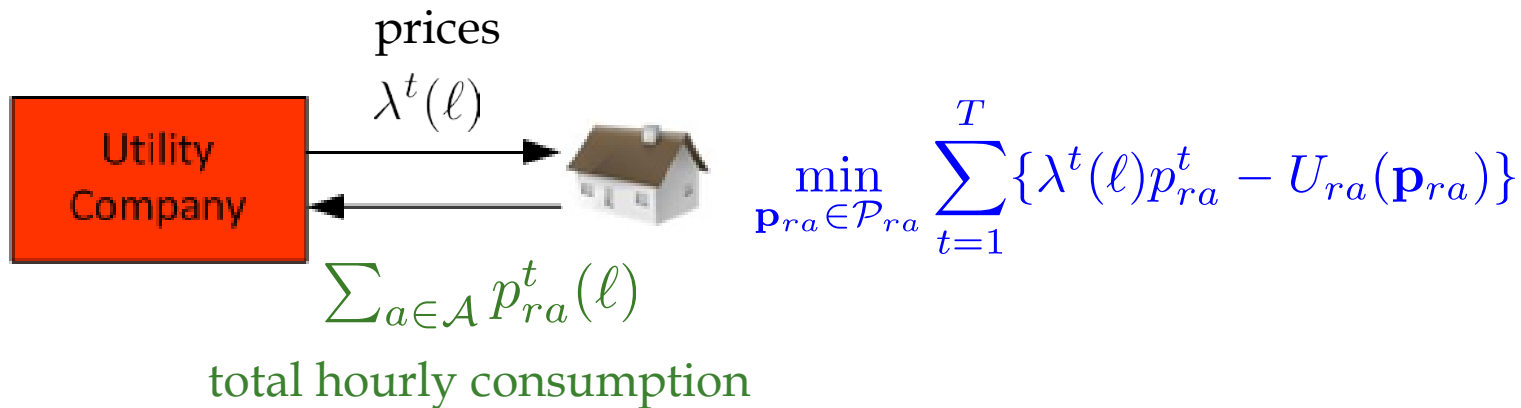
Demand-supply balance

- Lagrange multiplier λ^t for supply-demand balance
- **Upshot**
 - Sub-problems for utility and smart meters are separated
 - Privacy respected

Distributed DR algorithm

- **Schedule update:** At the utility company and smart meters

$$\min_{0 \leq s^t \leq s^{\max}} \{C^t(s^t) - \lambda^t(\ell)s^t\}$$

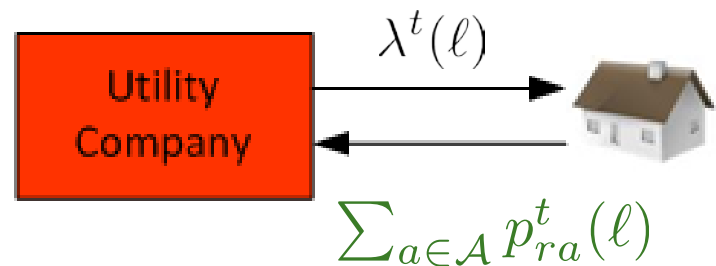


- **Multiplier update:** At utility company

$$\lambda^t(\ell + 1) = \left[\lambda^t(\ell) + \beta \left(\sum_{r=1}^R \sum_{a \in \mathcal{A}} p_{ra}^t(\ell) - s^t(\ell) \right) \right]^+$$

Lost AMI messages

- Messages in both ways may be lost
 - Not transmitted, due to failure
 - Not received, due to noise
 - Cyber-attacks



- Use the latest message available
- Convergence established for different lost-message patterns
 - Asynchronous subgradient method
- **Benefit:** **Resilience** to communication network outages



Plug-in Electric Vehicles

Plug-in electric vehicles

- PEVs feature batteries that can be plugged in
 - At end-user premises
 - At charging stations
- Benefits of high PEV penetration
 - Environmental: reduce carbon emissions
 - Economic: reduce dependency on oil
- Charging coordination is well motivated to avoid
 - Overloading of distribution networks [Clement-Nyns et al'10]
 - Creating new peaks



Charging coordination

- Fleet of vehicles $n = 1, \dots, N$ to charge on top of baseload D^t
- Fraction of charge (rate) r_n^t per slot $t = 1, \dots, T$
 - Vehicle plugged in at different slots $\Rightarrow 0 \leq r_n^t \leq r_n^{t,max}$
- Centralized charging coordination

$$\begin{aligned} \min_{\{\mathbf{r}_n\}} \quad & \sum_{t=1}^T C^t \left(D^t + \sum_{n=1}^N r_n^t \right) \\ \text{subj.to} \quad & 0 \leq r_n^t \leq r_n^{t,max} \\ & \sum_{t=1}^T r_n^t = R_n \end{aligned} \quad \left. \vphantom{\sum_{t=1}^T} \right\} \mathcal{E}_n \quad n = 1, \dots, N$$

- Convex and differentiable $C^t(\cdot) : \mathbb{R} \rightarrow \mathbb{R}$, e.g., $C^t(\cdot) := (\cdot)^2$

Distributed PEV scheduling

- Relies on Frank-Wolfe (FW) optimization method

- Identical per-vehicle partial gradients of the costs

$$\nabla_{\mathbf{r}_n} \left(\sum_{t=1}^T C^t \left(D^t + \sum_{n=1}^N r_n^t \right) \right) = \mathbf{g}, \quad n = 1, \dots, N$$

- At iteration i , vehicle n solves a linear program

$$\hat{\mathbf{r}}_n(i) \in \arg \min_{\hat{\mathbf{r}}_n \in \mathcal{E}_n} \hat{\mathbf{r}}_n^\top \mathbf{g}(i)$$

↑
auxiliary variable

Solution: charge first small entries of $\mathbf{g}(i)$

- Vehicle n updates $\mathbf{r}_n(i+1) = (1 - \eta(i)) \mathbf{r}_n(i) + \eta(i) \hat{\mathbf{r}}_n(i)$

$$\eta(i) := \frac{2}{i+2} \text{ or chosen via line-search}$$

Asynchronous updates

- To deal with *process delays* of EV controllers and/or *commutation failures*

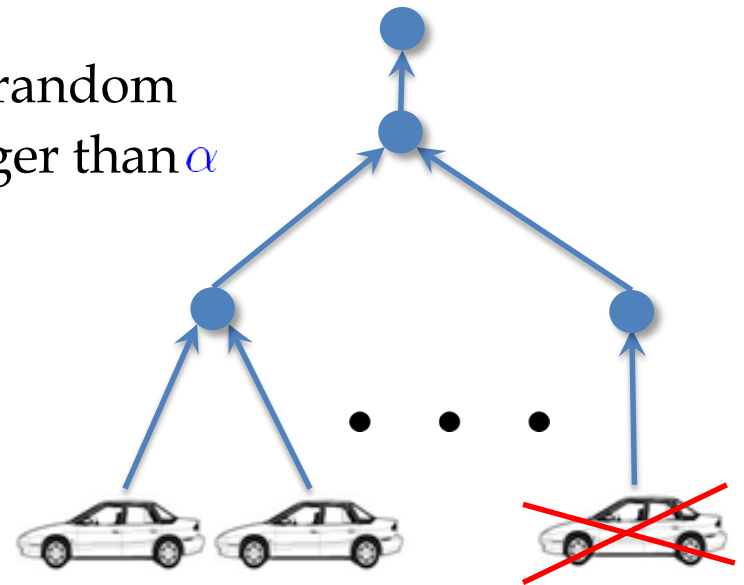
(as1) Lost updates occur independently at random

(as2) Probability of a successful update larger than α

- Guaranteed convergence with

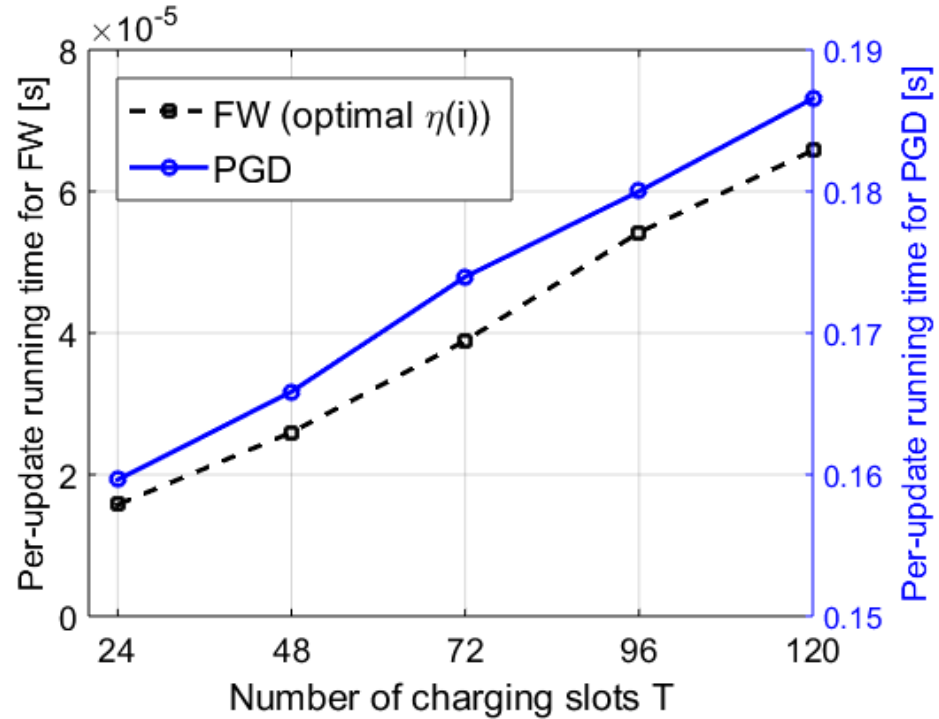
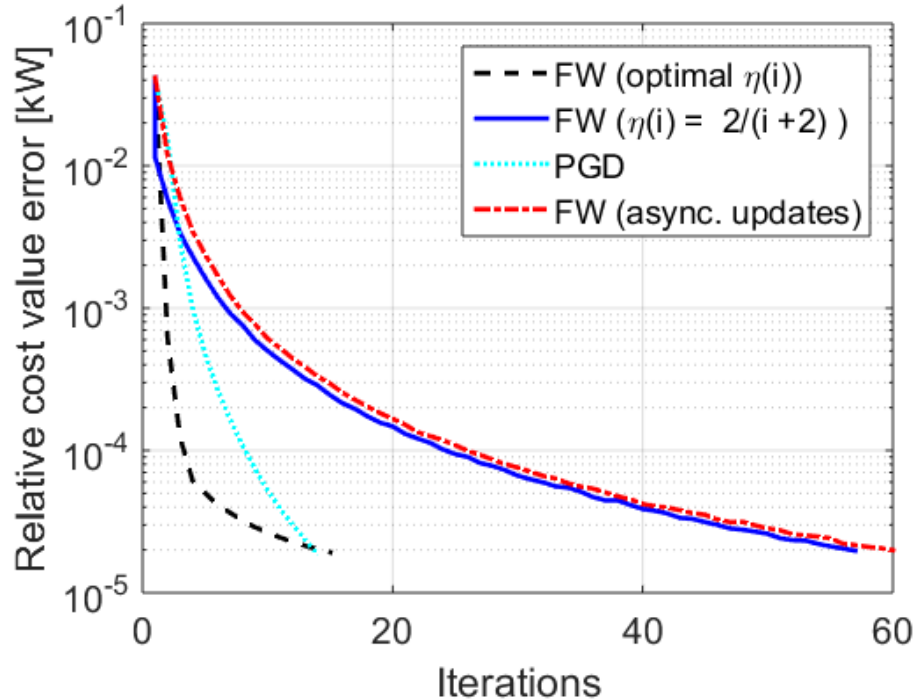
$$\eta(i) := \frac{2}{\alpha i + 2}$$

- $\mathcal{O}(1/i)$ convergence rate in expectation



Numerical tests

- 51 out of 52 EVs are updated in an asynchronous setting



- Projected gradient descent (GD) and ADMM must project (expensive!)
- Speed-up advantage* of FW thanks to simple updates

Take-home messages

- Distributed and robust PSSE
 - Non-convexity tackled via semidefinite relaxation
 - Decentralized estimation via *ADMM*
 - Sparse outlier models for robustness to “bad data”
- Distributed OPF
 - Semidefinite relaxation is tight for radial microgrids
 - *ADMM* solver for decentralized multiphase OPF
- Distributed DR
 - Decentralized management through *dual decomposition*
 - Resilience to lost AMI messages
- Distributed EV charging
 - Scalable and decentralized scheduler via *Frank-Wolfe* iteration
 - Robust to random communication outages

Thank you!